

Guillermo Escobar P cod 256243

Exercises 6.5-3

Write pseudocode for the procedures HEAP-MINIMUM, HEAP-EXTRACT-MIN, HEAP-DECREASE-KEY, and MIN-HEAP-INSERT that implement a min-priority queue with a min-heap.

Solucion :

El pseudocódigo es :

Heap-Minimum(A)

return A[1]

Heap-Extract-Min(A)

if heap-size[A] < 1

    then error "heap underflow"

min  $\leftarrow$  A[1]

A[1]  $\leftarrow$  A[heap-size[A]]

A[heap-size]  $\leftarrow$  A[heap-size[A]-1]

Min-heapify(A,1)

return min

Heap-Decrease-Key(A,i,key)

if A[i] < key

    then error "new key is greater than current key"

A[i]  $\leftarrow$  key

while i > 1 and A[i] < Parent(i)

    do exchange A[i]  $\longleftrightarrow$  A[Parent(i)]

    i  $\leftarrow$  Parent(i)

Min-Heap-Insert(A,key)

heap-size[A]  $\leftarrow$  heap-size[A] + 1

A[heap-size[A]]  $\leftarrow -\infty$

Heap-Decrease-Key(A,heap-size[A],key)

Exercises 6.5-4

Why do we bother setting the key of the inserted node to  $-\infty$  in line 2 of MAX-HEAP-INSERT when the next thing we do is increase its key to the desired value?

Solucion :

Since line 1 of the `Heap-Increase-Key` procedure needs a value for  $A[i]$ , the one to be inserted. Moreover, line 2 of the same procedure requires that such a value has to be smaller than the one to be increased to. Thus, we simply set it to be the smallest possible one.

#### Exercises 6.5-5

Argue the correctness of `HEAP-INCREASE-KEY` using the following loop invariant:

- At the start of each iteration of the **while** loop of lines 4-6, the array  $A[1 \dots \text{heap-size}[A]]$  satisfies the max-heap property, except that there may be one violation:  $A[i]$  may be larger than  $A[\text{PARENT}(i)]$ .

#### Exercises 6.5-6

Show how to implement a first-in, first-out queue with a priority queue. Show how to implement a stack with a priority queue. (Queues and stacks are defined in [Section 10.1](#).)

#### Exercises 6.5-7

The operation `HEAP-DELETE( $A, i$ )` deletes the item in node  $i$  from heap  $A$ . Give an implementation of `HEAP-DELETE` that runs in  $O(\lg n)$  time for an  $n$ -element max-heap.

#### Exercises 6.5-8

Give an  $O(n \lg k)$ -time algorithm to merge  $k$  sorted lists into one sorted list, where  $n$  is the total number of elements in all the input lists. (*Hint*: Use a min-heap for  $k$ -way merging.)

#### Problems 6-1: Building a heap using insertion

The procedure `BUILD-MAX-HEAP` in [Section 6.3](#) can be implemented by repeatedly using `MAX-HEAP-INSERT` to insert the elements into the heap. Consider the following implementation:

```
BUILD-MAX-HEAP'(A)
1  heap-size[A] ← 1
2  for i ← 2 to length[A]
3      do MAX-HEAP-INSERT(A, A[i])
```

- a. Do the procedures BUILD-MAX-HEAP and BUILD-MAX-HEAP' always create the same heap when run on the same input array? Prove that they do, or provide a counterexample.
- b. Show that in the worst case, BUILD-MAX-HEAP' requires  $\Theta(n \lg n)$  time to build an  $n$ -element heap.